# Exploiting Live Virtual Machine Migration

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**Black Hat DC** 





## . Introduction to VM migration

- . Live migration security
- Exploiting live migration
- . Future attacks and wrap-up



• Transfer of a VM from one physical machine to another with little or no service downtime



High Availability Enhanced Mobility

#### **Dynamic Load Balancing**

### **Live Migration Methodology**

- . Minimize service downtime
- Minimize migration duration
- Migration Types:
  - . Stop-and-copy (S-C)
  - . Demand-migration (D-M)
  - . Iterative precopy (I-P)





- Stop source VM
- . Copy all pages over the network
- Start destination VM

## Stop and Copy

**Longest Service Downtime** 

**Shortest Migration Duration** 



- . Copy over critical OS structures
- Start destination VM
- . Page faults trigger network copy

## **Demand Migration**

**Shortest Service Downtime** 

**Longest Migration Duration** 



- . Iteratively copy pages over network
- . Keep copying dirtied pages until threshold
- At threshold, stop source VM, copy remaining pages, start destination VM
   Iterative Precopy

**Balances Service Downtime and Migration Duration** 

Method used by VMware/Xen





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#### **A Trip Down Memory Lane**



#### Physical machines

- . Machine state protected by MMU/hardware
- . Physical attacks (firewire device DMA)
- Virtual Machines
  - . VM state protected by VMM/hypervisor
  - . Software attacks (weak VMM isolation)

Can we break any more isolation boundaries?



#### Of course! Functionality always usurps security!

- Migration-enabled VMs
  - . Full VM state exposed to network
  - Trades off security for management capabilities
  - Authentication, confidentiality, isolation concerns





- . Migration data plane
  - . Network transit path over which migration occurs
- . Security of data plane
  - . Unauthenticated, insecure migration data plane
- . Full access granted to VM state
  - . OS/kernel memory
  - Application state
  - . Sensitive data, passwords, keys, etc
- VMware and Xen migrations vulnerable



- . Breach of data plane means game over
  - Entire virtual machine may be compromised
  - . Kernel, userspace applications, data
- . Requirement for breach
  - Manipulate traffic along migration path between source and destination VMM
- Need to perform MITM attack
  - ARP/DHCP spoofing
  - . DNS spoofing/poisoning
  - . IP/route hijacking



#### Passive Attacks

- Snarf sensitive data, passwords, keys in memory
- . Active Attacks
  - . Manipulate authentication services
    - sshd, /bin/login, pam, etc
  - Manipulate kernel structures
    - slip rootkits into memory







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#### . Xensploit

- Non-weapons-grade proof-of-concept tool
- . Works against Xen and VMware migrations
- Attack classes
  - . VM application/userland exploits
  - OS/kernel exploits
  - VMM subversion



- sshd authentication bypass
  - . Identify pubkey authentication routines
  - Manipulate to allow unrestricted root access
  - . Access wide-open after migration completes
- Cron daemon shellcode injection
  - · Privileged, inconspicuous daemon
  - . Inject HTTP GET + execve shellcode
  - Payload fetch/exec on next find\_jobs invocation



## sshd authentication bypass

- Before migration:
  - attacker denied access to VM
- During migration
  - Xensploit manipulates the in-memory object code of sshd as it crosses the wire
- After migration
  - attacker achieves unrestricted root access to VM



Attacker attempts to gain root access to the target virtual machine via ssh

**jonojono@jonojono ~ \$** ssh root@vm-test Password: Password: Password: Permission denied (keyboard-interactive). **jonojono@jonojono ~ \$** 

. Attacker is denied access to the VM

#### sshd Authentication Code

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Source code from OpenSSH's auth2-pubkey.c:

```
if (key != NULL)
                key free(key);
        xfree(pkalg);
        xfree(pkblob);
#ifdef HAVE CYGWEN
        if (check nt auth(0, authctxt->pw) = 0)
                authenticated = 0;
#endif
        return authenticated;
/* return 1 if user allows given key */
static int
user key allowed2(struct passwd *pw, Key *key, char *file)
        char line[SSH MAX PUBKEY BYTES];
        int found key = \Theta;
        FILE *f;
        u long linenum = 0;
        struct stat st;
        Key *found;
        char *fp;
        /* Temporarily use the user's uid. */
        temporarily use uid(pw);
        debug("trying public key file %s", file);
        /* Fail quietly if file does not exist */
        if (stat(file, &st) < 0) {
                                                                               196.2-9
                                                                                             63°
```

#### **During Migration**



• Xensploit manipulates the object code of sshd's authentication routines as it crosses the wire

805da77:	Of 84 23 fd ff ff	je	805d7a0 ⊲user_key_allowed2+0x80>
805da7d:	89 3c 24	mov	%edi,(%esp)
805da80:	e8 37 e5 fe ff	call	804bfbc <fclose@plt></fclose@plt>
805da85:	8d 85 8c df ff ff	lea	Oxffffdf8c(%ebp),%eax
805da8b:	89 44 24 04	mov	%eax,Ox4(%esp)
805da8f :	c7 04 24 15 0e 08 0	8 movl	\$0x8080e15,(%esp)
805da96:	e8 d5 28 01 00	call	8070370 <logit></logit>
805da9b:	e8 20 bd 01 00	call	80797c0 <restore_uid></restore_uid>
805daa0:	81 c4 9c 20 00 00	add	\$0x209c,%esp
805daa6:	31 c0	xor	%eax,%eax
805daa8:	Slo	рор	%ebx
805daa9:	5e	рор	%esi
805daaa :	5f	рор	%edi
805daab :	5d	рор	%ebp
805daac :	c3	ret	
805daad	8d 76 00	lea	OxO(%esi),%esi
0805dab0 ⊲user_key_allowed>:			
805dab0:	55	push	%ebp
805dab1:	89 e5	mov	%esp,%ebp

 Xensploit injects a *mov \$0x1,%eax* instruction into user\_key\_allowed2, returning 1 (true)

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Attacker again attempts to gain root access
 via ssh on the target virtual machine

jonojono@jonojono ~ \$ ssh root@vm-test
Last login: Thu Oct 18 15:18:37 2007 from jonojono.eecs.umich.edu
fjbox1 ~ # id
uid=0(root) gid=0(root) groups=0(root),1(bin),2(daemon),3(sys),4(adm),6(disk),10(wheel),11(flop
py),20(dialout),25(at),26(tape),27(video),1006(vmware)
fjbox1 ~ #

- No authentication is necessary as sshd's routines have been manipulated by Xensploit
- . Root access is granted to the attacker



- . Kernel manipulation
  - . Direct access to in-memory kernel image
  - . More complexity but more power
  - . Leverage all your DMA attack payloads
- . Stealthy backdoor drop
  - network/syscall/ioctl trigger
- Full-blown VMBR hoisting



- . Virtual Machine-Based Rootkits
  - . Slip in extra virt layer a la SubVirt/Blue Pill/Vitriol





- Mangle migration payload
  - Exploit a vulnerability and subvert VMM
- . Leverage Xen dom0 vulns
  - Present in Xen daemon migration routines
  - . <= 3.1.0 release vulnerable</p>
  - . Undoubtedly more...
- . Instantly own all hosted VMs
  - And all future migrated VMs!

#### **Subverting the VMM**



. Xen's libxc/xc\_domain\_restore.c:

```
unsigned long region_pfn_type[MAX_BATCH_SIZE];
. . .
   for (;;)
       int j, nr mfns = 0;
. . . .
       if ( !read_exact(io_fd, &j, sizeof(int)) )
. . .
       if(j = -1)
. . .
       if(j = -2)
       if ( j = 0)
. . .
       if ( j > MAX BATCH SIZE )
. . .
       if ( !read exact(io fd, region pfn type, j*sizeof(unsigned long)) )
```

- No check for signed integer j < 0
- Stack overflow of region\_pfn\_type in Xen VMM

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Lots more juice in the migration orange!

- Fraudulent migration requests
  - Owned VMMs snarfing up VMs
- . False resource advertising
  - Migration-enabled load balancing
- . Future attacks inevitable
  - Increased functionality
  - Increased complexity
  - . Increased security risk



- Encryption goes a long way!
- Fingerprinting migrations
  - . Reconnaissance / targeting
  - . Enabled by iterative-precopy method
  - . Similar to VBR attacks
- Increased complexity
  - . Full PKI adds considerable deployment complexity
- Not currently implemented!

#### **Vendor Response**



#### . VMware

- . Use separate network for migration paths
- . Use hardware-based crypto cards
- VMotion/Virtual Infrastructure 3 vulnerable
- . XenSource
  - . Consult vendor/distribution for security fixes
  - . Latest open-source release still at risk
  - Unsure of migration status in XenServer4
- Microsoft Hyper-V
  - Will they get it right?



- . VM migration paradigm
  - . VERY useful functionality
  - . Awareness of security risk necessary
  - Better isolation, access control, authentication
- . Until then...
  - Severe weaknesses exist in extensively deployed systems
  - . Valuable weapon for pentester/attacker





# **QUESTIONS?**

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